COSC 311: ALGORITHMS MINI 3

Solutions

1. Quicksort. Here's an unsorted array. Run quicksort on the array, showing what the array looks like *after* each partition step is complete, assuming you are using the last element as the partition element (you should *not* show the intermediate steps taken while partitioning). You can show the left and right recursions in the same picture.

9	11	6	15	2	10	4	7	16	13	12	5	8	3	14	1	
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Solution:

Bold and underlined numbers indicate those positions which are fixed on each level of the recursion.

First partition around 9:



Now partition around 5 (on the left) and 13 (on the right):

Now partition around 2 (left-most section), 8 (second section), 10 (third section), and 15 (right-most section):

Now all of the "singletons" are at the base case, and we do one more partition in the sections with two numbers—partition around 4, 7, and 12:

And now all remaining sections are at the base case, and we're done.

2. Induction. Suppose you were going to prove by induction that a certain claim holds for all powers of 2, starting with 1. What steps would you take to prove this? Explain in your own words why those steps form a complete, convincing proof.

Solution:

a. Base Case: I will show the claim holds for $n_0 = 1$. b. Induction Step: Assuming the claim holds for $n = 2^k$ (Induction Hypothesis), I'll show that the claim then holds for $n = 2^{k+1}$.

Explanation: Having proven base case, I showed the claim holds for $n=2^0$.

Then I know from the proof of my inductive case that the claim will also hold for $n = 2^{(0+1)} = 2^1$. Applying the inductive step on $n = 2^1$ again, I know the claim holds for $n = 2^{1+1} = 2^2$. and so on

The same logic can be applied to show that the claim holds for any positive integer powers of 2. Note that I do not have to show directly that the claim holds for 2, 4, 8, and so on. The inductive step takes care of all of these cases at once!